Collective Sequential Pattern Mining in Distributed Evolving Data Streams

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Abstract. The advances in processing and communication techniques resulted in a multitude of emerging applications that interact with streams of data. Traditional data mining systems store arriving data, collect them for later mining, and make multiple passes over the collected data. Unfortunately, these systems are prohibitively slow when they deal with data streams with massive amounts of data arriving at high rates. This paper introduces a new model for mining sequential patterns on distributed data streams environments. It focuses on evolving data streams that originate from multiple distributed sources. Moreover, the mining process is achieved without compromising the privacy of the individual data streams of the participant nodes. Simulation results show that the proposed model scales linearly with the number of distributed nodes. In addition, it reduces the overhead in the distributed mining process.

Keywords: Sequential Pattern Mining, Collective Data Mining, Distributed Data Streams

1. Introduction

Mining sequential patterns in data streams has attracted significant attentions [1, 2, 3]. However, several needs of real-world applications that exhibit distributed nature of data streams have not been met. Several applications utilize data streams originating from multiple distributed sources such as online intrusion detection systems, sensor networks, and web-click streams. These applications are geographically distributed stream-oriented systems by their nature. Additionally, data from all sources should be gathered in a single location to reach the result. In several situations, the cost of centralizing data can be prohibitive since streaming data rates may exceed the capabilities of the storage, communication, and processing infrastructure. In addition, privacy issues may be encountered in such environments. Consequently, additional challenges may be imposed on the design and development of mining techniques to be able to interact with these huge continuous multiple data streams. Thus, there is a real need for a new model of mining sequential patterns in distributed data streams that preserves the privacy of the individual nodes. It mainly helps in the situations where participant nodes may not fully trust each other in terms of the distribution of their own data.

Sequential pattern mining in distributed databases has been addressed by several researchers [4, 5, 6, 7]. In addition, several sequential pattern-mining algorithms in data streams are introduced in [8, 9, and 10]. However, these approaches are not capable of dealing with distributed streams. Mainly because collecting data streams in a single location before processing them has a prohibitive cost and can compromise the owner's privacy. On the other hand, the problem of mining distributed data streams was addressed in many areas of data mining such as mining association rules [11, 12], clustering [13, 14], and mining frequent itemsets [15, 16]. Conversely, mining sequential patterns in multiple data streams is addressed in [17] and [18]. However, the MILE algorithm introduced in [17] is a one-time fashioned algorithm that cannot retain previous mining results. Hence, it may take prohibitive amount of time in re-mining. This problem has been

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avoided in IAspam algorithm introduced in [18] that incrementally mines across-streams sequential patterns to maintain the newest mining results. However, it utilizes a centralized setting strategy, which relies on a sliding window at the same working node to read and sample data streams. Meanwhile, a framework for Collective Data Mining (CDM) has been proposed in [19]. The main objective of CDM is to ensure that partial models produced from local data at the different sites are correct. Moreover, these models can be utilized as building blocks for forming a global model. To the best of our knowledge, no study has been achieved to mine sequential patterns from distributed data streams in a distributed manner, which is the major motivation behind this work. Instead of gathering all data streams from different sources on a single location and performing data mining, a new model based on collective data mining is presented in this paper. This model relies on multiple working nodes, which mine data streams locally and forward only the summarized information to a coordinator node for a global mining process.

The rest of the paper is organized as follows: next section presents the problem formulation followed by the details of the proposed model. Then, a theoretical analysis of the proposed model is introduced followed by a detailed simulation study. Finally, the paper is summarized in the last section.

2. Problem Formulation

Given a set of \( d \geq 1 \) data streams \( DS^1, DS^2 \ldots DS^d \) on \( d \) identical nodes \( N_1, N_2 \ldots N_d \) and each stream \( DS^i \) contains a stream of sequences \( DS^i = s^i_1, s^i_2 \ldots \) etc. Let \( DS \) be the sequence-preserving union of these streams \( DS = DS^1 \cup DS^2 \cup \ldots \cup DS^d \). If \( \alpha \) is a sequential pattern then \( \alpha \) is a local frequent sequential pattern at node \( i \) if and only if the support of \( \alpha \) (the ratio of the number of sequences containing \( \alpha \) to the number of all sequences) in the local data stream \( DS^i \) is greater than or equal to the local minimum support \( \sigma_l \). In addition, \( \alpha \) will be a global frequent sequential pattern if and only if the support of \( \alpha \) in \( DS \) is greater than or equal to the global minimum support \( \sigma_G \).

The problem of collaborative sequential pattern mining is to find all sequential patterns \( GS \), whose support value is equal to or greater than a fixed global minimum support threshold \( \sigma_G \) for the known part of \( DS \) at a given time. Furthermore, the problem of privacy-preserving sequential pattern mining in distributed data streams is to discover sequential patterns embedded in \( DS \) by imposing a major constraint on the mining process. This constraint is to prevent any sharing of private data among participant nodes during the mining process.

3. Proposed Model

The proposed model tries to minimize the communication requirements among the nodes. In addition, it does not need raw data exchange among participant nodes. All nodes are assumed identical and one of the nodes is selected randomly to act as a coordinator node. The coordinator node collects stream summary from the rest of the nodes. The proposed model accomplishes the mining process without compromising the privacy of individual data streams at different participant nodes.

The existence of a large number of participant nodes and the rapid incoming streams may lead to a significant bottleneck at the coordinator node. This is mainly because all participant nodes should periodically send some information to the coordinator node. Hence, this may result in excessive communication and high space requirements at the coordinator node for processing a large number of incoming sequences. The load on the coordinator node could be significantly reduced by arranging the nodes in a hierarchical communication structure. At any node in the intermediate level \( i \), a set of frequent sequential patterns combined with their supports is received from the nodes at the lower level \( (i - 1) \). Then, integration is done at the intermediate node and a new set of frequent sequential patterns is mined with the level minimum support \( \sigma_i \). This process continues until the root node is reached and a set of global sequential patterns is mined. Consequently, the proposed model consists of three phases; mining sequential patterns at local nodes, an integration phase, and the global mining phase as shown in Fig. 1. On the other hand, Fig. 2 shows pseudo-code of each phase of the proposed model.

3.1. Mining Sequential Patterns at Local Nodes
SPEDS introduced by us in [20] is adopted to mine sequential patterns at local nodes. It uses a lexicographical tree $T$ to hold the subsequences extracted from the data stream so far. Each path from the root of the tree to a node represents a sequence where the items are the nodes along the path. Each node in the tree contains a tilted-time window table constructed based on a logarithmic time scale. The entries of this table represent the numbers of sequences seen in the data stream in the corresponding tilted-time window.

At any participant node $N_i$ corresponding to a data stream $DS_i$, a lexicographical tree $T_i$ is constructed and SPEDS is applied locally to mine the set of local frequent sequences $LS_i$. It utilizes a local minimum support threshold $\sigma_i$, and a significance threshold $\varepsilon$ ($0 \leq \varepsilon \leq \sigma_i$). After processing a predefined number of batches $n$ of local sequences at a participant node $i$, a set of local frequent sequences $LS_i$ combined with a list $C_i$ of their corresponding support are mined and forwarded to the coordinator node for further global mining.

The support of a sequence $s$ at any given time using the time fading scheme described in [20] is calculated by Eq. (1).

$$
|s| = \sum_{f_k > 0} \omega^{n-k} f_k + \sum_{f_k=0} \omega^{n-k} \alpha L_k .
$$

where the fading factor $\omega < 1$, $n$ is the number of all received batches, $f_1, f_2, \ldots, f_n$ are the numbers of occurrences of $s$ in tilted-time windows $t_1, t_2, \ldots, t_n$ respectively, $\alpha$ is the batch support threshold ($0 \leq \alpha \leq \varepsilon$), and $L_k$ is the number of all sequences in the $k^{th}$ tilted-time window.

A sequence $s$ is frequent if its support $|s| \geq (\sigma - \varepsilon) N$, where $\sigma$ is the minimum support threshold, $\varepsilon$ is the significance threshold ($0 \leq \varepsilon < \sigma$), and $N$ the number of sequences computed by:

$$
N = \sum_{k=1} \omega^{n-k} L_k .
$$

Fig. 1. The architecture of the proposed model
Mining local frequent sequential patterns at local node N

Repeat

While the number of the received batches < n Do
   Apply SPEDS on DS
   For every sequence s in T_coordinator
      Calculate |s| using Eq. (1)
      If |s| ≥ (\sigma - \varepsilon)N_L Then
         add s to LS and |s| to C
   Send LS, combined with a list C to the coordinator node
   Initialize tree T_coordinator
End Repeat

Integration phase at the coordinator node

For any incoming LS and C at time T, Do
   Repeat
      For every sequence s with support \sigma in LS and C
         If not end of pruning period Then
            If s is not in T_coordinator
               Insert s in T_coordinator
               Set the corresponding count to zero
            Else
               Increment count of the corresponding node by \sigma
            End if
            Scan T_coordinator
            For each node in T_coordinator
               Update tilted-time table
               Set count to zero
            End for
            Prune tree
         Else
            For each node in T_coordinator
               Set count to zero
            End for
            Prune tree
         End if
      Until end of pruning period
   End Repeat

Mining global frequent sequential patterns at the coordinator node

Calculate N_L using Eq. (4)
   For every sequence s in T_coordinator
      Calculate |s| using Eq. (3)
      If |s| > (\sigma_L - \varepsilon)N_L Then
         add s to GS
End for

Fig. 2. The pseudo-code of the three phases of the proposed model

Integration Phase. At the coordinator node, a lexicographical tree T_coordinator similar to the one at local nodes is created and upon receiving LS_i and C_i from participant node i, an integration is started by inserting all incoming local frequent sequences into T_coordinator. SPEDS algorithm is modified and applied at the coordinator node to mine the set of global frequent sequences GS. Eq. (1) above is modified at the coordinator node to be:

\[ |s| = \sum_{1 \leq k \leq m, f_k > 0} \omega^{m-k} f_k + \sum_{1 \leq k \leq m, f_k = 0} \omega^{m-k} (\alpha_L - \varepsilon)N_L. \]  

(3)

where \( m \) is the number of all received sequences and \( N_L \) is calculated using a modified version of Eq. (2), which is given by Eq. (4).

\[ N_G = \sum_{k=1}^{m} \omega^{m-k} N_L. \]  

(4)

The pruning period \( \delta \) is represented by the number of incoming sequences, and it is a multiple of the window size. Every time the pruning period is up, the tree is pruned by calculating \(|s|\) for each node using Eq. (3) and if \(|s| \leq \varepsilon N_G\) then the entire sub-tree rooted at this node is pruned.

Mining Global Sequential Patterns at the Coordinator Node. At the coordinator node, upon requesting the set of global frequent sequences GS, the time fading scheme presented in [20] is applied on T_coordinator using the global minimum support \( \sigma_G \) where \( \sigma_G > \sigma_L \).

4. Theoretical Analysis

The local mining process at the participant nodes is based on SPEDS, which has been proven to produce no false negatives [20]. Assume there is a sequence \( s' \) that is frequent and was not included into the global frequent sequences. Two cases for \( s' \) can be identified:
1. $s'$ was a frequent sequence at the local nodes $N_{i_1}, N_{i_2} \ldots N_{i_n}$ with frequencies $f_{i_1}, f_{i_2} \ldots f_{i_n}$ in the tilted-time windows $t_1, t_2 \ldots t_3$ at the coordinator node respectively. Hence, $f_{i_1}, f_{i_2} \ldots f_{i_n} \geq (\sigma_L - \epsilon) N_L$, where $\sigma_L$ is the local minimum support and $N_L$ is calculated using Eq. (2).

2. $s'$ was not frequent in the remaining local nodes $N_{j_1}, N_{j_2} \ldots N_{j_m}$ with frequencies $f_{j_1}, f_{j_2} \ldots f_{j_n}$ in tilted-time windows $j_1, j_2 \ldots j_3$ at the coordinator node respectively. Hence, $f_{j_1}, f_{j_2} \ldots f_{j_n} < (\sigma_L - \epsilon) N_L$ and the corresponding counts in entries $j_1, j_2 \ldots j_3$ of its tilted-time window table would be zeros. However, during calculating the sequence support using Eq. (3), its frequency is assumed to be $(\sigma_L - \epsilon) N_L$, which is greater than the true count of the sequence at the corresponding nodes. If such a sequence is not pruned from the tree then its support calculated using Eq. (3) is greater than the true count of the sequence. In this case, if it was not generated as an output then it would not be frequent. If such a sequence was pruned $n$ times from the tree then every time it is pruned from the tree $|s'| \leq e N_G$, where $e$ is the significance threshold $(0 \leq e < \sigma_G)$ and $|s'|$ is its support calculated by Eq. (3). Consequently, if it was not an output, it would not be frequent. Hence, sequence $s'$ does not exist, which leads to the following claim:

**Claim 1.** The proposed model produces no false negatives.

SPEDS guarantees that all false positives have true support of at least $(\sigma - \epsilon) N$. Where $N$ is the stream length, $\sigma$ is the minimum support threshold, and $\epsilon$ is the significance threshold $(0 \leq \epsilon < \sigma)$. To calculate the minimum support of false positives in the proposed model for mining sequential patterns from distributed data streams, let us first calculate the maximum error $e_{\text{max}}$ in the support of a sequence using Eq. (3). Assume that there is a sequence $s'$ that has the minimum support at the coordinator node, for such a sequence to be an output from a local node and forwarded to the coordinator node, its minimum local support at the local node(s) should be $(\sigma_L - \epsilon) N_L$. Hence, the minimum global support $|s'|_{\min}$ for any sequence in $T_{\text{coordinator}}$ using Eq. (3) is $|s'|_{\min} = \sum_{1 \leq k \leq \text{min}} \omega^{\text{min}} (\sigma_L - \epsilon) N_L$ and from Eq. (4): $|s'|_{\min} = (\sigma_L - \epsilon) N_G$. This leads to $e_{\text{max}} \leq (\sigma_L - \epsilon) N_G$. Consequently, for any sequence $s'$ that has true support $c$ with Eq. (3) utilized to calculate $|s'|$ then:

$$|s'| \leq c + (\sigma_L - \epsilon) N_G.$$  \hspace{1cm} (5)

If $s'$ was included in the output then $|s'| \geq (\sigma_G - \epsilon) N_G$, which is the minimum support for a sequence to be an output. On the other hand, from Eq. (5): $|s'| \leq c + (\sigma_L - \epsilon) N_G$, hence, for any sequence $s'$ to be an output it should have true count $c$ satisfying $c + (\sigma_L - \epsilon) N_G \geq (\sigma_G - \epsilon) N_G$. Consequently, it leads to the following claim:

**Claim 2.** In the proposed model, all false positives have true support of at least $(\sigma_G - \sigma_L) N_G$.

5. **Experimental Results**

Several experiments have been carried out with a fading factor $\omega = 0.5$, a batch size of 100 sequences, $\alpha = 0.0095$, and $\epsilon = 0.01$. Fig. 3 shows the results of this experimental study.
Synthetic data streams are adopted in the experiments. They have been generated using the sequential data generator introduced in *IlliMine* system package. The first experiment is carried out to show the effect of changing local minimum support $\sigma_L$ on the size of transmitted data from any distributed node. The number of processed batches at any distributed node is 1000 batches, $\sigma_L$ is varied and the average size of transmitted data to the coordinator node in bytes is measured.

The results of this experiment are shown in Fig. 3a. As shown in this figure, if $\sigma_L$ is high then there will be only a limited number of sequential patterns and, hence, the pattern length will be short.

The second experiment is carried out to assess the communication load at the coordinator node. In this case, the number of distributed nodes is changed and the communication load at the coordinator node is measured in terms of the size of received data from all distributed nodes. The number of processed batches at any distributed node is 1000 batches, $\sigma_L = 0.14$. As shown in Fig. 3b the communication load scales linearly with the number of distributed nodes. These results are compared with the communication load at the coordinator node if the mining process is performed in a centralized setting (Fig. 3c). As shown in this figure, there is a considerable reduction of the data transmitted to the coordinator node when adopting the
proposed model. This is mainly because the coordinator node only collects the stream summary from the participant nodes. This reduction in the transmitted data contributes to the scalability of the proposed model.

The third experiment is carried out to analyze the timing requirements at the coordinator node for the integration and mining of the global frequent sequences and generating the results. The number of processed batches at any distributed node is 1000 batches, $\sigma_L = 0.14$ and $\sigma_G = 0.2$. The number of distributed nodes is changed during the experiment. Then, the execution time of the integration phase and time required to generate the results at the coordinator node are measured. As shown in Fig. 3d, the average execution time for outputting the global frequent sequences scales linearly with the number of distributed nodes, while Fig. 3e shows that the execution time of integrating local frequent sequences at the coordinator node increases with the number of distributed nodes.

6. Summary

Most prior work on sequential pattern mining in data streams considers the case of a single stream or multiple data streams in a centralized setting. Hence, a distributed data stream mining model has been proposed in this paper that is able to mine sequential patterns from multiple distributed evolving data streams. It is proven that the proposed model produces no false negatives and imposes a lower bound of the support of false positives. Simulation study has been carried out to analyze the performance of the proposed model. Simulation results show that the proposed model reduces the communication overhead in the distributed mining process. Most importantly, it scales linearly with the number of distributed nodes, which contributes to the scalability of the proposed model.

7. References


